

Cheat Sheet 1

Proof General

Ctrl-C	Ctrl-Enter	Ctrl-X	Ctrl-F
	Goto the cursor		Open a file
Ctrl-X	Ctrl-S	Ctrl-X	Ctrl-C
	Save		Quit

Basic notations

Logic	Prop
\top, \perp	True, False
$\neg p$	$\sim p$
$p \wedge q$	$p \wedge q$
$p \vee q$	$p \vee q$
$a = b$	$a = b$
$a \neq b$	$a <> b$
$p \Rightarrow q$	$p \rightarrow q$
$p \Leftrightarrow q$	$p \leftrightarrow q$
$\forall x \in A. \forall y. q(x, y)$	<code>forall (x:A) y, q x y</code>
$\exists x \in A. p(x)$	<code>exists x:A, p x</code>

Paper	Coq
Lemma 1. <i>For all natural numbers n and m if $m \leq n$ then the following equation holds:</i> $n - m + m = n$ <i>Proof.</i> Trivial □	Theorem good_name : <code>forall n m,</code> $m \leq n \rightarrow$ $n - m + m = n.$ Proof. (* your proof *) Qed.

Basic commands

Require Import ssreflect ssrbool.
Load libraries ssreflect and ssrbool

Variable name : type.
Declares a variable

Theorem name : statement.
State a theorem

Proof.
Start the proof of a theorem

Qed.
Terminate the proof of a theorem

Terminology

Context $\left\{ \begin{array}{l} P : \text{nat} \rightarrow \text{Prop} \\ x : \text{nat} \\ h : P \ x \end{array} \right.$

The bar -----

Goal $\left\{ \underbrace{\text{forall } y, y = x \rightarrow P \ y}_{\text{Stack}} \underbrace{\phantom{\text{forall } y, y = x \rightarrow P \ y}}_{\text{Conclusion}} \right.$

Top is the head of the stack, here y

Basic proof commands

done.

Prove the goal by trivial means, or fail

0 <= n \rightarrow

move=> x px.

Introduce x and P x naming them x and px

forall x, \rightarrow $\begin{array}{l} x : T \\ px : P \ x \end{array}$
P x \rightarrow Q x \rightarrow G -----
Q x \rightarrow G

(**move=>** x _ . to throw away P x completely)

move: x px.

Generalise x and px

$\begin{array}{l} x : T \\ px : P \ x \end{array}$ -----
 ----- \rightarrow forall x,
Q x \rightarrow G P x \rightarrow Q x \rightarrow G

(**move:** x (px) . to leave a copy of px in the context)

apply: H.

Apply H to the current goal

H : A \rightarrow B -----
 ----- \rightarrow A
B

case: ab.

Eliminate the conjunction, disjunction or absurd

ab : A \wedge B -----
G \rightarrow -----
A \rightarrow B \rightarrow G

ab : A \vee B ----- \rightarrow ----- -----
G A \rightarrow G B \rightarrow G

ab : False -----
G \rightarrow

case: exg3.

Eliminate the existential quantification

exg3 : **exists** n, 3 < n -----
 ----- \rightarrow forall n,
G 3 < n \rightarrow G

case: x.

Perform a case analysis on x

x : nat ----- -----
 ----- \rightarrow P 0 forall x,
P x P (S x)

elim: x.

Perform an induction on x

x : nat ----- ----- -----
 ----- \rightarrow P 0 forall x,
P x P x \rightarrow P (S x)

split.

Prove a conjunction

----- ----- -----
A \wedge B \rightarrow A B

left.

Prove a disjunction choosing the left part.

----- -----
A \vee B \rightarrow A

right.

Prove a disjunction choosing the right part.

$$\frac{}{A \vee B} \rightarrow \frac{}{B}$$

exists n.

Prove an existence giving a witness.

$$\frac{n : \text{nat}}{\text{exists } x, P \ x} \rightarrow \frac{n : \text{nat}}{P \ n}$$

rewrite Eab.

Rewrite with Eab left to right

$$\frac{\text{Eab} : a = b}{P \ a} \rightarrow \frac{\text{Eab} : a = b}{P \ b}$$

rewrite -Eab.

Rewrite with Eab right to left

$$\frac{\text{Eab} : a = b}{P \ b} \rightarrow \frac{\text{Eab} : a = b}{P \ a}$$

have pa : P a.

Open a new goal for P a. Once resolved introduce a new entry in the context for it named pa

$$\frac{a : T}{G} \rightarrow \frac{a : T}{P \ a} \quad \frac{a : T}{\text{pa} : P \ a} \rightarrow \frac{}{G}$$

suffices pa : P a.

Open a new goal with an extra item in the context for P a named pa. When resolved, it asks to prove P a

$$\frac{a : T}{G} \rightarrow \frac{a : T}{\text{pa} : P \ a} \rightarrow \frac{a : T}{P \ a}$$

Notations for natural numbers: nat

"n.+1 :: " := (S n)

"n.-1 :: " := (predn n)

m + n := (addn m n)

m * n := (muln m n)

m <= n := (leq m n)