YicesQS 2024, an extension of Yices for quantified satisfiability

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YicesQS is a solver derived from Yices 2, an open-source SMT solver developed and distributed by SRI International. It extends Yices to supports quantifiers for complete theories, and is unrelated to the support of quantifiers in Yices for UF. Its core algorithm is a generalization of counterexample-guided quantifier instantiation (CEGQI) [Dut15] that can be seen as a form of lazy quantifier elimination. YicesQS submits quantifier-free queries to one of Yices's core solvers CDCL(T) or MCSAT [dMJ13]. It is written in OCaml and uses the OCaml bindings for the Yices C API; it indirectly relies on the libpoly library for arithmetic. In the 2024 SMT competition, YicesQS entered logics BV, NRA, NIA, LRA, and LIA (single-query, non-incremental tracks), as in 2023 and 2022. The exact commits are as follows:

https://github.com/disteph/yicesQS	commit 4661444
https://github.com/SRI-CSL/yices2_ocaml_bindings	commit 38318a6
https://github.com/SRI-CSL/yices2	commit 5082c90
https://github.com/SRI-CSL/libpoly_ocaml_bindings	commit 8e9dc78
https://github.com/SRI-CSL/libpoly.	commit 7a4dedc

and the exact command is

```
main.exe -under 20 -cdclT-mcsat 700 $1
```

Algorithmic approach. The main algorithm in YicesQS is QSMA [BGLV23]. It does not modify the structure of quantifiers in formulas: it does not prenexify formulas and, more importantly, it does not skolemize them to avoid introducing uninterpreted function symbols. This departs from the standard architecture for quantifier support consisting of keeping a set of universally quantified clauses, to be grounded and sent to a core SMT-solver for ground clauses. Instead, YicesQS sees a formula as a 2-player game, in the tradition of Bjørner & Janota's *Playing with Quantified Satisfaction* [BJ15] and earlier work on QBF. YicesQS's algorithm is designed to answer queries of the following form:

"Given a formula $A(\overline{z},\overline{x})$ and a model $\mathfrak{M}_{\overline{z}}$ on $\overline{z}$, produce either
- SAT($U(\overline{z})$), with $U(\overline{z})$ under-approx. of $\exists\overline{x}\ A(\overline{z},\overline{x})$ satisfied by $\mathfrak{M}_{\overline{z}}$; or
- UNSAT($O(\overline{z})$), with $O(\overline{z})$ over-approx. of $\exists\overline{x}\ A(\overline{z},\overline{x})$ not satisfied by $\mathfrak{M}_{\overline{z}}$;

where under-approximations and over-approximations are quantifier-free."

To answer such queries, YicesQS calls Yices's feature *satisfiability modulo a model*, while the production of under- and over-approximations leverages *model interpolation* and *model generalization*. When the input formula is in the exists-forall fragment, the algorithm degenerates to the CEGQI one used in Yices' $\exists\forall$ solver [Dut15] using quantifier-free solving and *model generalization*. *Model*

interpolation, a form of which is used within Yices's MCSAT solver to solve quantifier-free problems [JD21], also becomes useful with more quantifier alternations than $\exists\forall$. It generalizes to non-Boolean inputs the notion of UNSAT cores, which has been used in the *quantified-problems-as-games* approach [BJ15].

Arithmetic (NRA, NIA, LRA, LIA). *Solving Modulo a Model*: We use Yices's MCSAT solver; the MCSAT approach is inspired by CDCL, building a tentative solution model on a trail where Boolean and theory assignments are decided, propagated, and backtracked upon. MCSAT can natively take a partial model as input [JD21]. *Model interpolation*: We use the model interpolants natively produced by Yices's MCSAT solver. A lemma that is learnt at decision level 0, defeating the input model, constitutes a model interpolant. Learnt lemmas arise from theory-specific mechanisms for explaining conflicts, which in the case of arithmetic is leveraging Cylindrical Algebraic Decomposition (CAD) [JdM12, Jov17]. *Model generalization*: We mainly use model-projection (based on CAD once again). For NRA, the presence of division in benchmarks departs from the theoretic applicability of YicesQS's algorithm for complete theories, because of division-by-zero (which also makes the theory undecidable). Yices's CAD-based model-projection in NRA does not support division. When YicesQS needs to perform model generalization with a formula involving division, it cannot use CAD model-projection and resorts to generalization-by-substitution, which is a generic mechanism already used in Yices's $\exists\forall$ solver [Dut15]. Resorting to generalization-by-substitution for NRA also means that YicesQS's algorithm may not terminate.

Bitvectors (BV). *Solving Modulo a Model*: In 2024 we use primarily Yices's CDCL(T) solver, which itself relies on bitblasting and, ultimately, Yices's internal SAT-solver. The partial input model is represented as equalities between bitvector variables and their values, and given to Yices as assumptions. *Model interpolation*: We simply use UNSAT cores produced by Yices's CDCL(T) solver. *Model generalization*: We use invertibility conditions from Niemetz et al. [NPR$^+$18], including ϵ-terms, in combination with generalization-by-substitution. For the BV theory, the cegqi solver from [NPR$^+$18] is probably the closest to YicesQS. Note: instead of the CDCL(T) solver, it is possible to use Yices's MCSAT solver, which support quantifier-free bitvectors [GLJD20].

Changes since 2023. We implemented a look-ahead optimisation related to section 5 of our QSMA paper [BGLV23]. As another optimisation that we implemented, the top-level call to QSMA does not need to produce an under- or over-approximation, only the recursive calls do, while the top-level call only needs to conclude sat or unsat. For BV, rather than using MCSAT as in the previous years, we now primarily use CDCL(T) (i.e. bitblasting), which produces presumably poorer interpolants but is faster. The MCSAT-based approach does solve benchmarks that the CDCL(T)-based approach does not, so after 700 seconds we restart the solving with MCSAT in order to solve some extra benchmarks. For arithmetic, YicesQS 2024 benefits from recent optimisations in the underlying MCSAT solver (e.g., extended propagation mechanism and decision heuristics).

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