

Consistency of Task-Based Authorization Constraints in Workflow System

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Introduction

- Workflow: consists a collection of tasks that are organized to facilitate some business process specification
 - Purchase order processing
 - Processing tax refunds

Introduction

- Security Challenge
 - Administration of security-relevant information
 - Specification of authorization policies and constraints
 - Solution
 - RBAC
 - Constraints

Introduction

- Constraints: restriction on the authorization
 - Separation of Duty
 - Binding of Duties
- Challenge for constraints in workflow system
 - Consistency

Goal

- Help workflow designers to define a sound constrained workflow authorization schema such that
 - each workflow instance can complete
 - each user or role authorized to perform a task will have an opportunity to perform that task in some workflow instance

Related Work

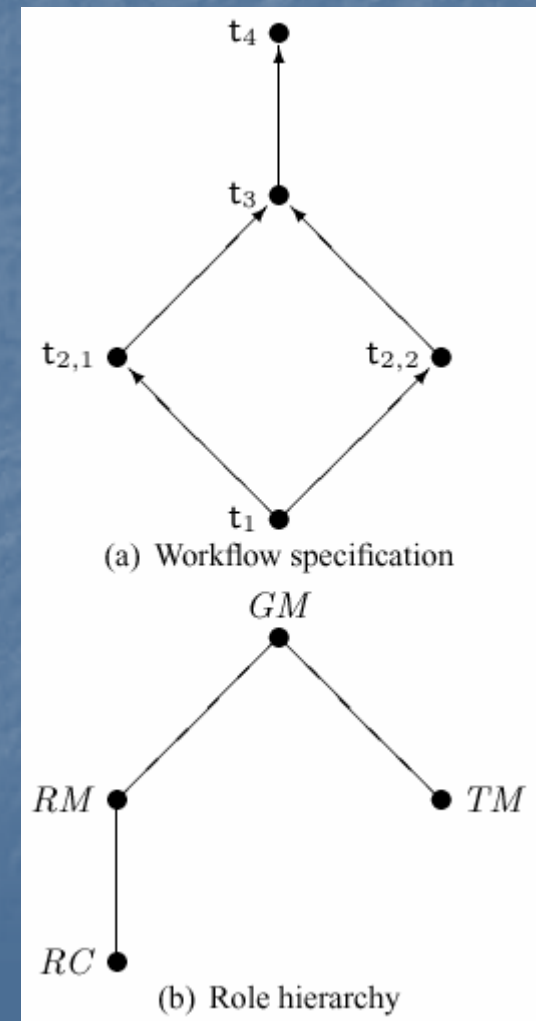
- E. Bertino, E. Ferrari, V. Atluri: TISS 99
 - Role planning and user planning
 - + Constraints are consistent
 - Complicated
- J. Wainer, P. Barthelmeß, A. Kumar: IJCIS 03
 - Override constraints when workflow can not be completed
 - + Completing workflow instance
 - Weaken security

Our Contribution

- Give explicit definition of constraint workflow authorization schema
- Develop simple specification of authorization constraints
- Define consistency rules which can be checked efficiently

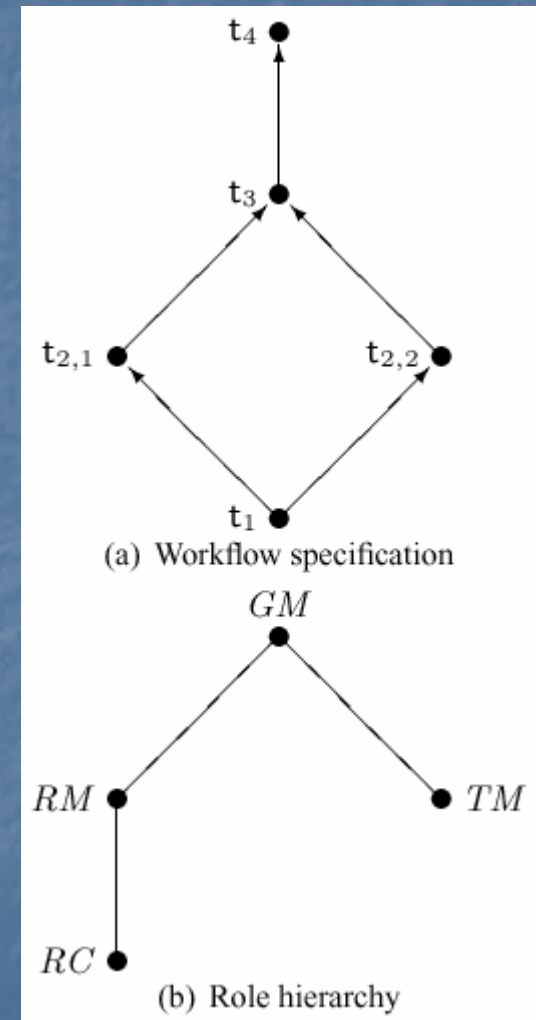
Workflow Example

- Producing payments for tax refunds
 - t_1 : a clerk prepares a check for a tax refund
 - t_2 : the check is approved or denied by two different managers
 - t_3 : A third manager makes a final decision based on the decisions made in t_2
 - t_4 : a clerk issues or voids the check based on the decision of t_3



Constraints

- c_1 : The two instances of t_2 should be performed by different users
- c_2 : t_2 and t_3 should be performed by different users
- c_3 : t_1 and t_4 should be performed by different users
- c_4 : t_2 must be performed by a role that is more senior than the role that performed t_1 unless t_1 and t_2 are performed by GM
- c_5 : t_1 and t_2 must be performed by different users
- c_6 : At least three roles should perform the workflow instance



Constraint Workflow Authorization Scheme

- $(T, \leq, \rho)$, where T is a (partially ordered) set of tasks and ρ is a function from T to the set of natural numbers indicating the number of occurrences of each task in the workflow
- Role hierarchy: $(R, \leq)$
- User-role association: $UA \subseteq U \times R$
- Task-role association: $PA \subseteq T \times R$
- Constraint-task association: $PC \subseteq C \times T$
- $R(t) = \{r' \in R : \exists (t, r) \in PA, r' \geq r\}$
- $U(t) = \{u \in U : (u, r) \in UA, r \in R(t)\}$
- Task instance: (t_i, u_i, r_i)
- Workflow instance: $[(t_1, u_1, r_1), \dots, (t_n, u_n, r_n)]$

Task	Role
t_1	<i>RC</i>
t_2	<i>RM</i>
t_3	<i>RM</i>
t_4	<i>RC</i>

(c) Task-role assignment relation PA

User	Role
Alice	<i>RC</i>
Bob	<i>RM</i>
Carol	<i>RM</i>
Dave	<i>RC</i>
Eve	<i>GM</i>
Fred	<i>TM</i>

(d) User-role assignment relation UA

Constraint-Task assignment PC

- $PC \subseteq C \times T$ where $C = \{c_1, c_2, c_3, c_4, c_5, c_6\}$
 - c_2 : t_2 and t_3 should be performed by different users
 - $(c_2, t_3) \in PC$
 - c_1 : The two instances of t_2 should be performed by different users
 - $(c_1, t_2) \in PC$
 - c_6 : At least three roles should perform the workflow instance
 - $(c_6, t_1) \in PC, (c_6, t_2) \in PC, (c_6, t_3) \in PC, (c_6, t_4) \in PC$

Constraints Specification

- Entailment constraints
 - Execution of t is constrained by the execution of t' where $t' < t$
 - **Entailed task** of $t' : t$
 - $(ur, t', E, pred)$
 - $ur \in \{u, r\}$
 - $E \subseteq U(t')$ if $ur=u$ or $E \subseteq R(t')$ if $ur=r$
 - $pred \in \{=, \neq, <, >, \leq, \geq\}$
 - $pred(u_{t'}, u_t) = \text{true}$ if $ur=u$ or $pred(r_{t'}, r_t) = \text{true}$ if $ur=r$
 - **c_2** : t_2 and t_3 should be performed by different users
 - $c_2 = (u, t_2, U(t_2), \neq)$

Cardinality Constraints

- Local cardinality constraints
 - Impose restrictions on the number of users to execute a task
 - c_1 : The two instances of t2 should be performed by different users
 - $c_1=(2,2)$
- Global cardinality constraints
 - Impose restrictions on the number of roles to execute a set of tasks
 - c_6 : At least three roles should perform the workflow instance
 - $c_6=(T,3)$

Consistency

- A constrained authorization schema $((T, \leq, \rho), PA, PC)$ is **sound** if
 - All $(c, t) \in PC$ is well-formed
 - A constraint-task pair is **well-formed** if
 - E.g., $(ur, t', E, pred)$
 - $ur = u$ or r
 - $t' < t$
 - $E \subseteq U(t')$ if $ur = u$ or $E \subseteq R(t')$ if $ur = r$
 - $pred \in \{=, \neq, <, >, \leq, \geq\}$
 - For all $u \in U(t)$ and all $r \in R(t)$, there is a successful workflow instance in which t is executed by u or r

Constraint Interaction

- Two constraint-task pairs **interplay** if
 - The constraints are assigned to the same task
 - E.g., $((2,2), t_2), ((r, t_1, R(t_1), <), t_2)$
 - The constraints are assigned to a task and its entailed task
 - E.g., $((u, t', U(t'), \neq), t)$ and $((2, 2), t')$ where t is an entailed task of t'
 - The constraints are assigned to a set of restricted tasks in T'
 - E.g., for global cardinality constraint $(T', n_r), ((T', n_r), t)$ and $((T', n_r), t')$ where $t, t' \in T'$

Interplaying Sets CT

- Affected tasks in CT
 - $((2,2), t_2), ((r, t_1, R(t_1), <), t_2): \{t_2\}$
 - $((u, t', U(t'), \neq), t), ((2, 2), t'): \{t\}$
- Non-interplay set SC
 - any (c,t) in SC does not belong to any CT or they are in some CT but t is not an affected task of that CT
- Any $(c,t) \in SC$ should be **self-consistent**
- Any CT should be **inter-play-consistent**

Interplay consistency

- Self-consistent or inter-play consistent
 - Consistency rules for different constraint-task pairs
 - E.g., for $c=(u, t', E, pred) \in SC$ and $(c,t) \in PC$, then for any $\alpha \in E$, $U(t|t', \alpha, c) \neq \emptyset$ and $\bigcup_{\alpha \in U(t')} U(t|t', \alpha, c) = U(t)$
 - Computations required
 - Set computation

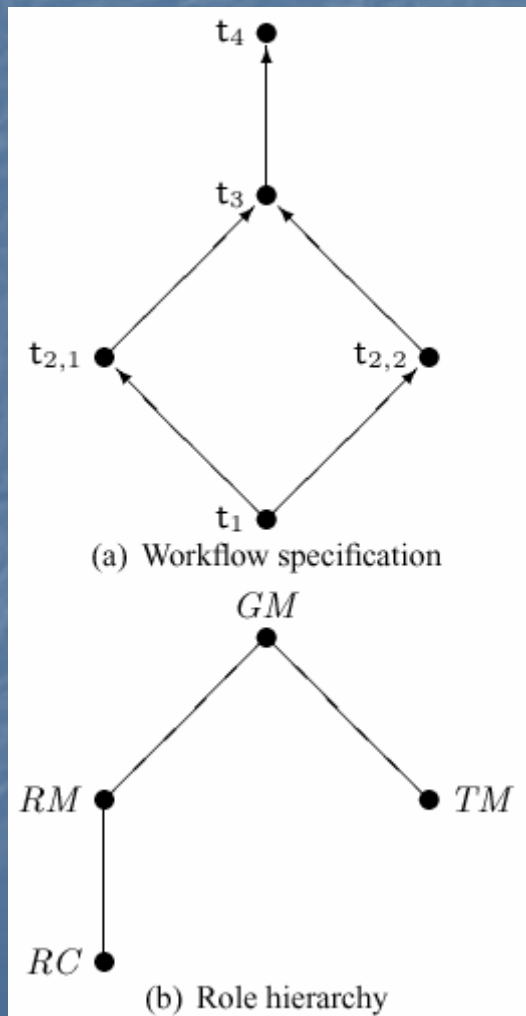
Main Result

- **Theorem:** a constrained workflow authorization schema $((T, \leq, \rho), PA, PC)$ is sound if and only if all constraint-task pairs in SC are self-consistent and all the inter-play sets are inter-play-consistent

Example

- Constraint Inconsistency of our example
 - E.g. $CT = \{(c_4, t_2), (c_1, t_2)\}$ inter-play
 - c_1 : the two instances of t_2 should be performed by different users
 - $c_1 = (2, 2)$
 - c_4 : t_2 must be performed by a role that is more senior than the role that performed t_1 unless t_1 and t_2 are performed by GM
 - $c'_4 = (r, t_1, R(t_1), <)$, $c''_4 = (r, t_1, GM, =)$
 - If Bob or Carol executes t_1 , then the workflow instance cannot complete at t_2

Example



Task	Role
t ₁	RC
t ₂	RM
t ₃	RM
t ₄	RC

(c) Task-role assignment relation PA

User	Role
Alice	RC
Bob	RM
Carol	RM
Dave	RC
Eve	GM
Fred	TM

(d) User-role assignment relation UA

Conclusion

- Give explicit definition of constraint workflow authorization schema
- Give a solution to constraints consistency problem in workflow system
 - Simple constraints specification
 - Efficient consistency rules checking