

# Formal Techniques for Analyzing Hybrid Systems

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## Motivation

We want to mathematically **model** a system and **analyze** its behavior using that model

- Systems such as cars, drones, airplanes that have both a **cyber** component (programs, protocols, controllers, planners, etc.) and a **physical** component

Most of modern science and engineering uses modeling and analysis

We, however, want to analyze things a bit more **formally**

There is plenty of **motivation** for why we want to model and analyze dynamical systems

- predict what will happen, increase confidence about behavior, establish correctness, reduce costs, etc.

## Dynamical Systems

A **dynamical system** consists of a notion of

- **state space**  $\mathbf{X}$  – defines the configurations of the system
- **collection of traces**  $F : \mathbf{T} \mapsto \mathbf{X}$  – defines all possible evolutions of the system, for some notion  $\mathbf{T}$  of time

The state space  $\mathbf{X}$  is usually described as the set of valuations of a given set of variables.

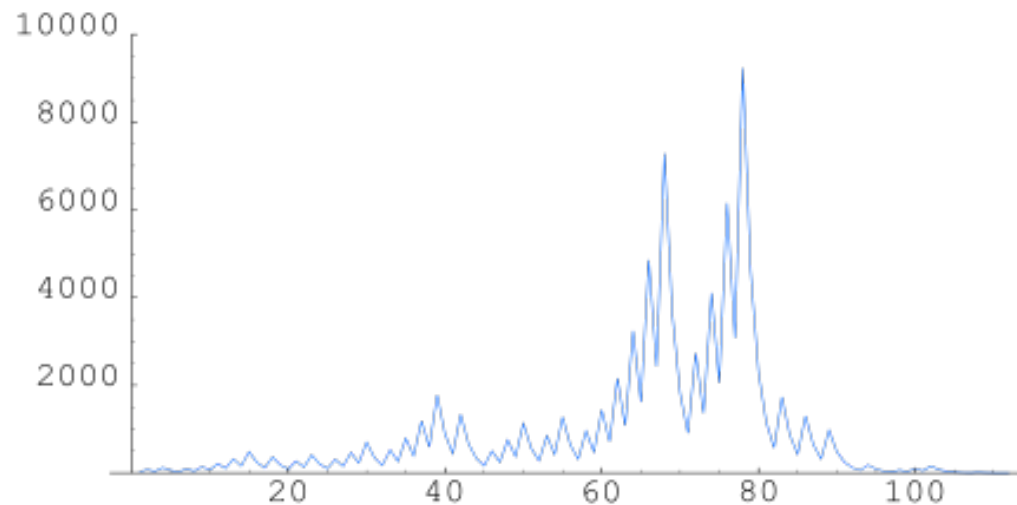
## Dynamical system: Example

**Example:** Consider one variable  $x$  taking values in the set of positive integers.

How to define the dynamics?

**Rule:**  $x(t + 1) = x(t)/2$  if  $x(t)$  is even, else  $3 * x(t) + 1$

A trace  $F(t)$  of the system when initially  $F(0) = 27$  is shown below.



## Dynamical Systems: Classifying based on $\mathbf{X}$ and $\mathbf{T}$

We can have different kinds of state space:

- discrete-space:  $\mathbf{X} \subset \mathbb{N}^n$
- continuous-space:  $\mathbf{X} \subset \mathbb{R}^n$
- hybrid-space:  $\mathbf{X} \subset \mathbb{N}^n \times \mathbb{R}^m$

We can have different kinds of “time”:

- discrete-time:  $\mathbf{T} = \mathbb{N}$
- continuous-time:  $\mathbf{T} = \mathbb{R}^{\geq 0}$
- hybrid-time:  $\mathbf{T} = \mathbb{R}^{\geq 0} \times \mathbb{N}$

The example above was a discrete-time discrete-space dynamical system

## Dynamical System: Exercise

1. Construct an example of a **discrete-time continuous-space** dynamical system

- Newton's method for finding a root of  $f$ :

$$x(t + 1) = x(t) - f(x(t))/f'(x(t))$$

2. Construct an example of a **continuous-time continuous-space** dynamical system

- Freely falling ball:

$$dy(t)/dt = v(t), \quad dv(t)/dt = -9.81$$

**Question:** What is the state space of these two systems?

## Dynamical System: Finite- and Infinite-State

A system is **finite-state** if there are finitely many elements in its state space  $X$ .  
It is **infinite-state** otherwise.

- All continuous-space systems are infinite-state.
- A dynamical system defined by a finite-state automaton is finite-state.

## Dynamical Systems: Classifying based on Dynamics

A system is **deterministic** if fixing  $F(0)$  fixes  $F$

Example: The  $3n + 1$  system was deterministic

In a **nondeterministic** system, the same state can have multiple “successors”.

Sometimes we assign a probability to each of the different options available at a state

This gives us a **stochastic** dynamical system.

Eg. Markov chain



## Dynamical Systems: Open and Closed

A system is **open** if its dynamics  $F$  is influenced by some input variables.

A system is **closed** if it has no input variables

The input variables can range over some predefined set of values

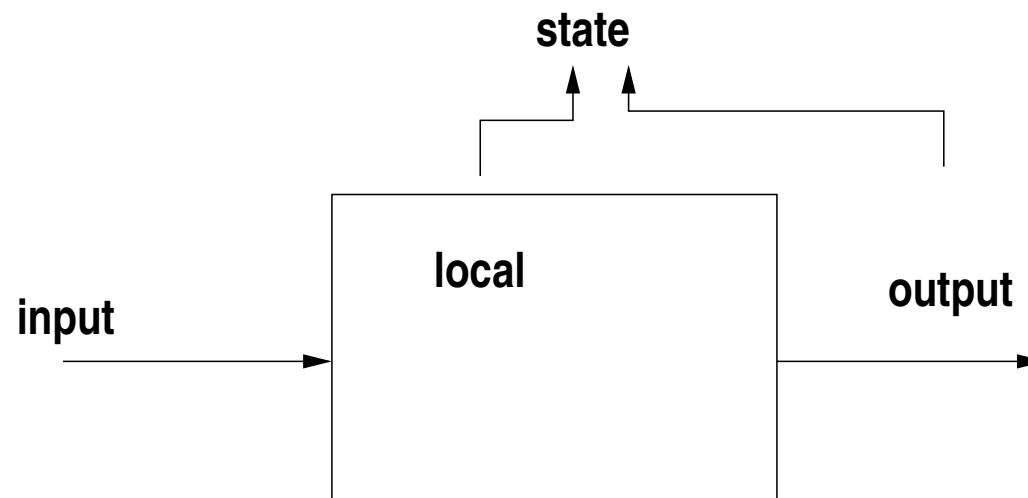
The input variables are **disjoint** from the state variables

**Example:** A human, or a controller, turns a heater on or off. The dynamics of the room temperature changes based on this **Boolean** input.

## Dynamical Systems: Observables

Some part of the state space of a system might be externally observable

This is modeled by defining a set of **output** variables, which is a **subset** of the state variables



## How to describe a dynamical system?

Recall the **two parts** of a dynamical system:

- **State space  $X$** : we can just list the variables and their types
  - we can identify local and output variables, and additional input variables if any
- Collection of **temporal traces**: We write rules that tell us how the system evolves **locally**
  - a program
  - a differential equation
  - a collection of rules
  - **compose** smaller dynamical systems

## Examples

Spring-mass system:

- State space:  $\mathbb{R}^2$ , valuations of  $x, v$
- Dynamics:  $\frac{dx}{dt} = v, \frac{dv}{dt} = -x$

A sorting program:

- State space:  $\mathbb{Z}^n$ , valuations of  $x_1, \dots, x_n$
- Dynamics: for any  $i, j$ : swap  $x_i, x_j$  if  $x_i > x_j$

**Exercise:** Classify these two dynamical systems. open/closed?  
det/non-det/stochastic? finite-state? discrete-time? discrete-space?

## Summary So Far

We defined dynamical systems, and the various classes of such systems

We saw a few examples

We discussed how to write a dynamical system

Next, we want to analyze dynamical systems to see if they have some desired property.

How do we specify the desired properties?

## Property Specification Language

Logic! To say things about temporal behavior, we have **temporal logics**.

$$\phi := \top \mid at \mid \neg\phi \mid \phi \vee \phi \mid \mathbb{F}\phi \mid \mathbb{G}\phi \mid \phi\mathbb{U}\phi$$

Here *at* denotes an atomic formula that can be evaluated on a **state**; e.g.  $x < 20$

What does a formula mean? It is a property of traces, and hence we should know if a given trace satisfies a given formula.

## Semantics for our Temporal Logic

$$F \models \top$$

$$F \models at \quad \text{if } at \text{ is true at state } F(0)$$

$$F \models \neg\phi \quad \text{if } F \not\models \phi$$

$$F \models \phi_1 \vee \phi_2 \quad \text{if } F \models \phi_1 \text{ or } F \models \phi_2$$

$$F \models \mathbb{F}\phi \quad \text{if } \exists t \geq 0 : F(t :) \models \phi$$

$$F \models \mathbb{G}\phi \quad \text{if } \forall t \geq 0 : F(t :) \models \phi$$

$$F \models \phi_1 \mathbb{U} \phi_2 \quad \text{if } \exists t \geq 0 : F(t :) \models \phi_2 \wedge \forall (t_1 < t) : F(t_1 :) \models \phi_1$$

The notation  $F(t :)$  denotes a suffix-trajectory of  $F$  that starts at state  $F(t)$

The time domain  $\mathbb{T}$  is assumed **unbounded**

## Temporal Properties

$\mathbb{F}$  operator says that eventually something happens

- $\mathbb{F}(\textit{good})$  : eventually something good happens

$\mathbb{G}$  operator says that always something holds

- $\mathbb{G}(\textit{notbad})$  : nothing bad ever happens

Relationship between temporal operators:

$$\mathbb{G}(\phi) = \neg\mathbb{F}(\neg\phi)$$

$$\mathbb{F}(\phi) = \top\mathbb{U}\phi$$



## Deadlock and Zeno

Recall: The time domain  $T$  is assumed **unbounded**

If a model **deadlocks**, then **time** ceases to progress

A hybrid system could have **zeno** behaviors.

Example: Bouncing ball

Reachable deadlock states, or zeno behaviors, are undesirable

Detecting if a system has such behavior is difficult

## Initial States

Before we can ask if a dynamical system satisfies a property, we need one last thing

We need to specify a set of initial states  $I \subseteq X$

We are only interested in trajectories  $F$  s.t.  $F(0) \in I$

## The Model Checking Problem

Given a dynamical system  $S = (\mathbf{X}, \mathbf{F}, I)$

- $\mathbf{X}$  - set of states
- $\mathbf{F}$  - collection of trajectories
- $I$  - set of initial states

And given a temporal property  $\phi$  (whose atomic formulas get evaluated on  $\mathbf{X}$ )  
determine if

$$F \models \phi$$

for every  $F \in \mathbf{F}$  s.t.  $F(0) \in I$

Notation:  $S \models \phi$

## Model Checking Problem: Example

Let  $S = (\mathbb{N}^{\{x\}}, \mathbf{F}, \mathbb{N}^{\{x\}})$  be the  $3n + 1$  system

An example of the model checking problem:

Is it the case that

$$S \models \mathbb{F}(x == 1)$$

This is an **open** problem

## Model Checking Problem: Example

For the same system  $S$ , consider:

- $S \models (x < 5) \Rightarrow \mathbb{G}(x < 15)$  ?
  - No
  - **Counter example:** The trace in  $\mathbf{F}$  where formula evaluates to false
  - 3, 10, 5, 16, 8, (4, 2, 1)\*
- $S \models (x < 5) \Rightarrow \mathbb{G}(x \leq 16)$  ?
  - Yes

## Summary

- Dynamical system:  $(\mathbf{X}, \mathbf{F}, I)$ 
  - each trajectory in  $F \in \mathbf{F}$  is a mapping from  $\mathbf{T}$  to  $\mathbf{X}$
- Property language - a temporal logic
  - a  $\phi$  is either true or false in any given trajectory
- Model checking problem:
  - Determine if  $S \models \phi$

## Tool Install

- Download and install SAL (executable)  
<http://sal.csl.sri.com>
- Install HybridSal relational abstractor:  
<http://www.csl.sri.com/users/tiwari/relational-abstraction/>