

Talk at IST Austria 9 March 2016, based on Invited talk, 12th International Conference on Distributed Computing and Internet Technology (ICDCIT), Bhubaneswar, India, January 2016

# Trustworthy Self-Integrating Systems

John Rushby

Computer Science Laboratory  
SRI International  
Menlo Park, California, USA

## Introduction: Systems of Systems

- We're familiar with systems built from components
- But increasingly, we see systems built from other systems
  - [Systems of Systems](#)
- The component systems have their own purpose
  - Maybe at odds with what we want from them
- And they generally have vastly more functionality than we require
  - Provides opportunities for unexpected behavior
  - Bugs, security exploits etc. (e.g., [CarShark](#))
- Difficult when [trustworthiness](#) required
  - May need to [wrap](#) or otherwise [restrict](#) behavior of component systems
  - And that means integration requires [bespoke engineering](#)

## Self-Integrating Systems

- But we can imagine systems that recognize each other and **spontaneously integrate**
  - Possibly under the direction of an “**integration app**”
  - Examples on next several slides
- Furthermore, separate systems often interact through shared “**plant**” **whether we want it or not** (**stigmergy**)
  - Separate medical devices attached to same patient
  - Car and roadside automation  
(autonomous driving and traffic lights)

And it would be best if they “**consciously**” integrated

- These systems need to “**self integrate**”
- And we want the resulting system to be trustworthy
- That’s a tall order

# Scenarios

- I'll describe some scenarios, mostly from medicine
- And most from [Dr. Julian Goldman](#) (Mass General)
  - “Operating Room of the Future” and
  - “Intensive Care Unit of the Future”
- There is [Medical Device Plug and Play \(MDPnP\)](#) that enables basic interaction between medical devices
- And the larger concept of “[Fog Computing](#)” to provide reliable, scaleable infrastructure for integration
- But I'm concerned with what the systems [do together](#) rather than the mechanics of their interaction

## Anesthesia and Laser

- Patient under general **anesthesia** is generally provided **enriched oxygen supply**
- Some throat surgeries use a **laser**
- In presence of **enriched oxygen**, laser causes **burning**, even **fire**
- Want laser and anesthesia machine to recognize each other
- **Laser requests reduced oxygen** from anesthesia machine
- But...
  - Need to be sure laser is talking to anesthesia machine connected to **this patient**
  - **Other** (or faulty) devices should not be able to do this
  - Laser should light only if oxygen **really is** reduced
  - In emergency, need to enrich oxygen should **override** laser

## Heart-Lung Machine and X-ray

- Very ill patients may be on a **heart-lung machine** while undergoing surgery
- Sometimes an **X-ray** is required during the procedure
- Surgeons **turn off** the heart-lung machine so the patient's **chest is still** while the X-ray is taken
- Must then remember to **turn it back on**
- Would like heart-lung and X-ray mc's to recognize each other
- X-ray requests heart-lung machine to **stop** for a while
  - **Other** (or faulty) devices should not be able to do this
  - Need a guarantee that the heart-lung **restarts**
- **Better**: heart lung machine informs X-ray of **nulls**

## Patient Controlled Analgesia and Pulse Oximeter

- Machine for Patient Controlled Analgesia (PCA) administers pain-killing drug on demand
  - Patient presses a button
  - Built-in (parameterized) model sets limit to prevent overdose
  - Limits are conservative, so may prevent adequate relief
- A Pulse Oximeter (PO) can be used as an overdose warning
- Would like PCA and PO to recognize each other
- PCA then uses PO data rather than built-in model
- But that supposes PCA design anticipated this
- Standard PCA might be enhanced by an app that manipulates its model thresholds based on PO data
- But...

## Patient Controlled Analgesia and Pulse Oximeter (ctd.)

- Need to be sure PCA and PO are connected to **same patient**
- Need to cope with **faults** in either system and in communications
  - E.g., if the app works by **blocking** button presses when an approaching overdose is indicated, then loss of communication could remove the safety function
  - If, on the other hand, it must **approve** each button press, then loss of communication may affect pain relief but not safety
  - In both cases, it is necessary to be sure that faults in the blocking or approval mechanism cannot generate **spurious button presses**
- This is **hazard analysis** and **mitigation** at integration time

## Blood Pressure and Bed Height

- Accurate **blood pressure sensors** can be inserted into intravenous (IV) fluid supply
- Reading needs correction for the **difference in height** between the sensor and the patient
- Sensor height can be standardized by the IV pole
- Some hospital beds have **height sensor**
  - Fairly **crude device** to assist nurses
- Can imagine an ICU where these data are available on the local network
- Then integrated by monitoring and alerting services
- But...

## Blood Pressure and Bed Height (ctd.)

- Need to be sure bed height and blood pressure readings are from **same patient**
- Needs to be an **ontology** that distinguishes height-corrected and uncorrected readings
- Noise- and fault-characteristics of bed height sensor mean that **alerts** should be driven from changes in **uncorrected reading**
- Or, since, bed height seldom changes, could synthesize a noise- and fault-masking **wrapper** for this value
- Again, **hazard analysis** and **mitigation** at integration time

## What's the Problem?

- Could build all these as **bespoke systems**
- More interesting is the idea that the component systems discover each other, and **self integrate** into a bigger system
- Initially will need an extra component, the **integration app** to specify what the purpose should be
- But later, could be more like the way human teams assemble to solve difficult problems
  - **Negotiation** on **goals**, exchange information on **capabilities**, **rules**, and **constraints**

## What's the Problem? (ctd. 1)

- Since they were not designed for it
- It's **unlikely** the systems **fit together perfectly**
- So will need **shims, wrappers, adapters** etc.
- So part of the problem is the “**self**” in self integration
- **How are these adaptations constructed during self integration?**

## What's the Problem? (ctd. 2)

- In many cases the resulting assembly needs to be **trustworthy**
  - **Preferably** do **what was wanted**
  - **Definitely** do **no harm**
- Even if self-integrated applications seem harmless at first, will often get used for critical purposes as users gain (misplaced) confidence
  - E.g., my Chromecast setup for viewing photos
  - Can imagine surgeons using something similar (they used Excel!)
- **So how do we ensure trustworthiness?**

## Aside: System Assurance

- State of the art in system assurance is the idea of a **safety case** (more generally, an **assurance case**)
  - An **argument** that specified **claims** are satisfied, based on **evidence** (e.g., tests, analyses) about the system
- System comes with **machine-processable online rendition** of its assurance case
  - Not standard yet, but Japanese DEOS project does it
  - Essentially a **proof**, built on premises **justified by evidence** (see my AAA15 paper, cf. one on Ontological Argument)
- **Ideally**: when systems self integrate, **assurance case for the overall system is constructed automatically from the cases of the component systems**
- Hard because **safety** often **does not compose**
  - E.g., because there are **new hazards**
  - Recall laser and anesthesia

## What's the Problem? (ctd. 3)

- While building the assurance case at self-integration time
- Likely must **eliminate** or **mitigate** some **hazards**
- May be able to do this by **wrappers**, or by **monitoring**
- Aside: **the power of monitors**
  - A monitor can be very simple
  - Can make a claim that it is **probably fault-free**
    - ★ **This** is the claim that verification delivers
  - Prob. of failure of **system** is then
    - ★ **prob. of failure** of operational component  
**times prob. monitor is fault-free**
  - Nb. **cannot** multiply probs. of failure
  - See TSE 2012 paper by Littlewood and me
- **How do these wrappers and monitors get built?**

## Models At Runtime (M@RT)

- If systems are to adapt to each other
- And wrappers and monitors are to be built at integration-time
- Then the systems need to know something about each other
- One way is to exchange models
  - Machine-processable (i.e., formal) description of some aspects of behavior, claims, assumptions
- This is Models at RunTime: M@RT
- When you add aspects of the assurance case, get Safety Models at RunTime: SM@RT (Trapp and Schneider)
- Most recent in a line of system integration concepts
  - Open Systems, Open Adaptive Systems, System Oriented Architecture

## Four Levels of SM@RT

- Due to Trapp and Schneider
- **Safety Certificates @ runtime** (feasible today)
  - Each system maintains its own local safety objective
  - But composed system may not be safe
- **Safety Cases @ runtime** (feasible tomorrow)
  - Component system safety cases guide adaptation
  - Integrated dynamically for safe & assured assembly
  - E.g., one system may need to demonstrate it delivers properties assumed by another
- **V&V @ runtime** (our goal, feasible soon)
  - May be that one system cannot deliver assumptions required by another
  - So adjustments needed
  - E.g., wrappers or monitors to exclude some class of faults
- **Hazard Analysis & Risk Assessm't at RT** (infeasible today)

## Example: SILF

### SILF: Semantic Interoperability Logical Framework

- Developed by NATO to enable dependable machine-to-machine information exchanges among Command and Control Systems
- Extensive ontology to describe content of messages exchanged
  - So in SM@RT terms, **ontological descriptions** (e.g., in **OWL**) are the models
- **Mediation mechanism** to translate messages as needed
  - **Synthesized at integration time**
- Mediation can be performed by centralized hub, or by wrappers at either the sender or receiver

## ONISTT and Onward

- **ONISTT** is an SRI project, **prototyped** ideas of SILF
  - Ad-hoc Prolog program synthesizes the mediator
    - ★ Now uses F-Logic and Flora2
  - Synthesis procedure can also decide when incompatibilities too great to meet purpose of integration
  - Used successfully to integrate live and virtual simulation systems for military training
- ONISTT achieves restricted form of **safety cases @ runtime**
- More general applications likely require **richer models than ontologies**
  - E.g., **state machines and formal specifications**
- How to perform synthesis on these?

## Synthesis as Exists/Forall Problem

- At integration time, systems need to synthesize wrappers, monitors, shims etc.
- Synthesis can be seen as a **generate and verify search problem**
  - Construct a candidate program
  - Try to formally verify that it meets specification
  - If not, generate new candidate and iterate
- Unrestricted search **will not work**
- Have human provide **template/sketch**, synthesis **fills in details**
- Simple example of a template for an invariant  $Ax + By = C$
- Formally, this can be expressed as

$$\exists A, B, C : \forall x, y : Ax + By = C \quad (1)$$

where  $x$  and  $y$  are program variables, and the parameters  $A$ ,  $B$ ,  $C$  must be instantiated by the synthesis procedure

- Note two-level quantification: **Exists/Forall (EF)**

## Synthesis as Exists/Forall Problem (ctd. 1)

- Variants on EF formulation can express
  - Invariant generation
  - Assumption synthesis
    - ★ Find the weakest environment in which a given component meets its requirements
  - Supervisory controller synthesis
    - ★ Design an algorithm to selectively disable component actions so that it satisfies some goal in the face of uncontrollable actions by the environment
  - Full synthesis
    - ★ Design an algorithm to achieve some goal
- So how do we solve EF problems?
- Start by solving **one-level problems**: Exists or Forall

## Synthesis as Exists/Forall Problem (ctd. 2)

- Satisfiability Modulo Theories (SMT)
- A breakthrough in automated theorem proving, 15 years ago
- Decides Boolean formulas over combination of theories
- ... Boolean formulas: e.g.,  $(x \leq y \vee y = 5) \wedge (x < 0 \vee y \leq x) \wedge x \neq y$   
... continued for many terms
- ... over combination of theories

e.g.,  $2 \times car(x) - 3 \times cdr(x) = f(cdr(x)) \supset$

$$f(cons(4 \times car(x) - 2 \times f(cdr(x)), y)) = f(cons(6 \times cdr(x), y))$$

Uses equality, uninterpreted functions, linear arithmetic, lists

- Can extend to one level of quantification  
(i.e., either Exists or Forall)
- There are many SMT solvers, honed by competition
- Routine to handle hundreds of thousands of terms in seconds

## Synthesis as Exists/Forall Problem (ctd. 3)

- **EF-SMT** solver uses an ordinary SMT solver as a **component**
  1. **Guess** (cleverly) instantiations for the **Exists** variables and **query** the SMT solver with the resulting **Forall** formula
  2. If this succeeds, we are done
  3. If it fails, use the result (i.e., **counterexample**) of the **Forall** query to help in finding the next instantiation of the **Exists** variables
- Key in making this efficient is to use (i.e., learn from) the result of failed verification (Forall) steps to prune the search space for subsequent synthesis (Exists) steps
- **Many SMT solvers being extended to EF solving** (e.g., **Yices**)

# Vision

- Systems **come together**
- **Exchange** models, assurance cases
- Under **guidance** of an **integration app**
  - Which expresses the **purpose** of the integration
    - ★ E.g., as a **template or sketch**
- **Connectors, wrappers, monitors, and shims** are **synthesized**
  - By **EF-SMT** solver
- And **system assurance case** is **composed** from those of **component systems**
- Delivers a **trustworthy integration**

## Conclusion

- Trustworthy self integration is within reach
  - For simple cases. . .
- Need theorem proving at integration time
  - To synthesize the connectors, monitors etc.
  - And to build the composed assurance case
- So a theorem prover will be at the heart of self integration
- In future, will likely also use learning to infer properties beyond supplied models
- Further ahead, will integrate highly autonomous systems
  - Numerous failures in HMI (e.g., Air France and Air Asia crashes) show this is difficult
- So must exchange more strategic information than SM@RT
- Maybe beliefs, desires, intent (BDI), even a system of ethics
- This is the future of IoT