

Composition of Critical Properties Lessons From Other Fields

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Composition

- Just a fancy word for **modular construction**
- We build systems from components
- Some of which are standard, some bespoke
- All tied together with “glue logic”
- We do this for economy, efficiency, quality
 - Good components are a **reusable asset**
- **Our reasoning about the system is based on what we know about the components and the way they are put together**
- What’s wrong with that?

Modern Automobiles

- Engine and its control system were traditionally developed separately from transmission and its control system
- But they actually have to work together
 - e.g., transmission tells engine to retard ignition while it changes gear
- Given the two components, it can take six months to integrate them
 - Not due to intended interactions such as the one above
 - But unanticipated consequences of interactions
 - And low-level details, like bus timing
- Delayed introduction of automated cruise, lane monitoring/following, automated parking, integration of powertrain with steering and suspension
- Reputations of some luxury manufacturers damaged

Composition Is Easy, No It's Hard

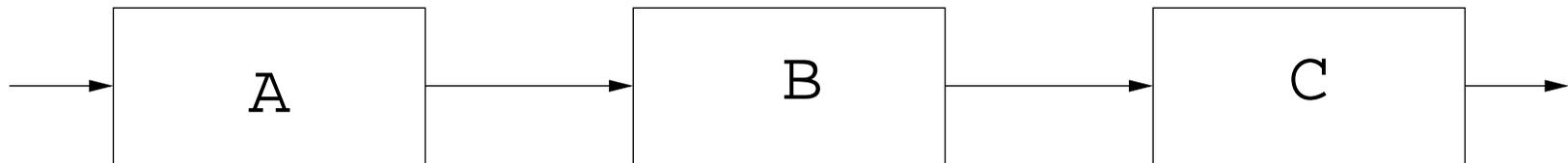
- It's easy when there is little interaction among components
 - e.g., 1960s automobiles, OB1
- It gets harder the more they interact
- Emergent behavior
 - Behavior of the system not found in any component
 - That's why we build systems—e.g., medical device PnP
- But then there's unanticipated, undesired emergent behavior
- In the limit, we get Normal Accidents (Perrow)
 - High interactive complexity
 - Tight coupling
- Challenge is to eliminate bad emergence, keep the good

System Properties

- Properties like safety, security, real-time guarantees
- These are properties of the **whole system**
 - e.g., where's the component that makes an airplane safe?
 - But a single component can easily make it unsafe
- i.e., these are **emergent properties**
- That's why the FAA certifies only airplanes and engines
 - OK, propellers too
 - Components certified only as part of an airplane or engine
 - Because you need to examine them in their context of interaction
- But this is becoming ruinously expensive, even infeasible
- Our goal is **compositional development** and **assurance**

It's About Interactions

- We must have **no unintended interactions**
- But in most systems almost every component can affect every other indirectly
- So we need to focus on the **direct** interactions



- And the nature (channels) of interaction
- And we need to be concerned about **un**intended consequences of intended interaction channels
- Particularly in the presence of **faults, malice**

Lessons From Other Fields

- The most sophisticated treatment is in embedded systems
 - They need real time, fault tolerance, safety
- They have the following concepts
 - Error-propagation boundaries
 - Elementary and composite interfaces
 - Fault-containment regions
 - Composability
- And from integrated modular avionics (IMA), we get
 - Partitioning
 - Determinism, time-triggered scheduling
- And from EU safety certification, we get
 - Argument-based assurance cases

Error-Propagation Boundaries

- Errors in a component should be detectable at its interface, **before** they propagate to other components
- Two kinds of errors: in **control** and in **data**
- Interfaces move data and use control (e.g., a protocol) to accomplish it
- Control errors are particularly destructive in real-time systems because they affect workload in the victim and hence ability to meet deadlines
- Control errors can be detected if there is redundant information
 - e.g., static common knowledge such as **fixed schedules**
 - Then have less or no need for data in control messages
 - ★ e.g., destination of message
- Consequences of control errors depend on interface

Elementary and Composite Control Interfaces

- **Composite** interfaces are those where control is **bidirectional**
 - Even when data flow is unidirectional
 - e.g., producer-consumer, queues

Problem is they allow errors to propagate in both directions

- **Elementary** interfaces have **unidirectional** control
 - Same direction as data flow
 - e.g., wait-free, lock-free, atomic registers
 - Such as Simpson's 4-Slot, Non-Blocking Write (NBW)

Errors propagate in only one direction

- Choice of data affects type of interface
 - **Event** vs. **state** messages
 - Events require confirmation: therefore composite

Fault-Containment Regions

- Two kinds of data errors
 - Send wrong **value**, send it at wrong **time**
- **Value errors** require redundancy and selection/voting
 - A host of delicate issues, understood by very few
 - ★ So you get **homespun designs**
 - ★ e.g., read incident/accident reports such as NTSB A07-65 through 86 (Predator), or A08-46, 47 (Eclipse)
 - A key idea is that of **fault containment region**
 - Required so that faults in redundant values will be **independent**
- **Timing errors** (e.g., babbling) are very destructive
 - Guaranteed elimination requires fixed schedules
 - e.g., static common knowledge, enforced by **bus guardians**
 - ★ Such as TTEthernet (used by Project Orion)

Composability

- aka. **preservation of prior properties**
- **A property established for a component or subsystem will not be invalidated by system integration**
- Even when other parts of the system have faults
- i.e., if components A and B do their thing within allocated processor and bus utilizations, rest of system must never invalidate this
 - Even when another component causes a processor exception
 - Or babbles on the bus
- **Composability is a stepping stone to compositionality**
- DO-297 talks of “**tiers of integration**”

Partitioning

- Allows components to **share** processor, bus resources
- By **eliminating unintended interactions**
- **Space partitioning**
 - Cannot read/write another component's memory
- **Time partitioning**
 - Cannot affect another component's access to processor/bus allocation
- Robust partitioning is key technology for integrated modular avionics (IMA)
- Now COTS technology from RTOS vendors
- And avionics buses: AFDX (weak), TTEthernet (strong)

Assurance and Certification

- We have **claims** or **goals** that we want to substantiate
 - Concerning some system property
- We produce **evidence** about the **product** and its development **process** to support the claims
 - E.g., **analysis** and **testing** of the product and its **design**
 - And **documentation** for the process of its **development**
- And we construct an **argument** that the evidence is **sufficient** to support the claims
- This is the intellectual basis for **all** certification regimes
- Claims and argument generally **implicit** in standards-based assurance, which focus on evidence to be produced
- Argument-based safety/security/dependability cases require **explicit** claims, evidence, argument

Compositional Assurance/Certification

- Assurance case may not decompose along architectural lines
 - Insight due to Ibrahim Habli & Tim Kelly
- Goes to the heart of what is an architecture
- A good one supports and enforces the assurance case
- We've now got enough background to see how to do this

Synthesis

We need:

- **Robust partitioning** to share processor and bus resources
- **Determinism** to control faults in the time domain
- **Redundancy** to tolerate faults in the data domain
- And **fault containment regions** so faults are independent
- **Elementary control interfaces** to provide error-propagation boundaries
- **Composability** so we can build things piecewise (layers)

Relationship to Security

- Adversary models for security are generally stronger than fault hypotheses for safety and fault tolerance
 - Active malice rather than Mother Nature
 - Though Mother Nature is assumed to be a strong cryptographer
 - ★ e.g., checksums and nuclear triggers
- Disclosure is more subtle than (most) faults
 - Any observable variation in behavior can be a side channel or covert channel that discloses sensitive information
- Manifestation of this is that security is not even a property

Properties and Security

- A **property** is a (possibly infinite) set of behaviors
 - **Safety property**: no bad thing happens
 - **Liveness property**: good things do happen (eventually)
 - **Any property** is the **intersection** of a safety property and a liveness property
- **The only things we can enforce are safety properties**
- Information flow security is not a property
 - It's a **hyperproperty**
 - Sets of sets of behaviors
- **But every hyperproperty can be enforced by a safety property**
 - e.g., information flow enforced by access control
 - May exclude some good behaviors
- We'll **enforce safety properties**, do **end-to-end analysis as hyperproperties**

The MILS Idea

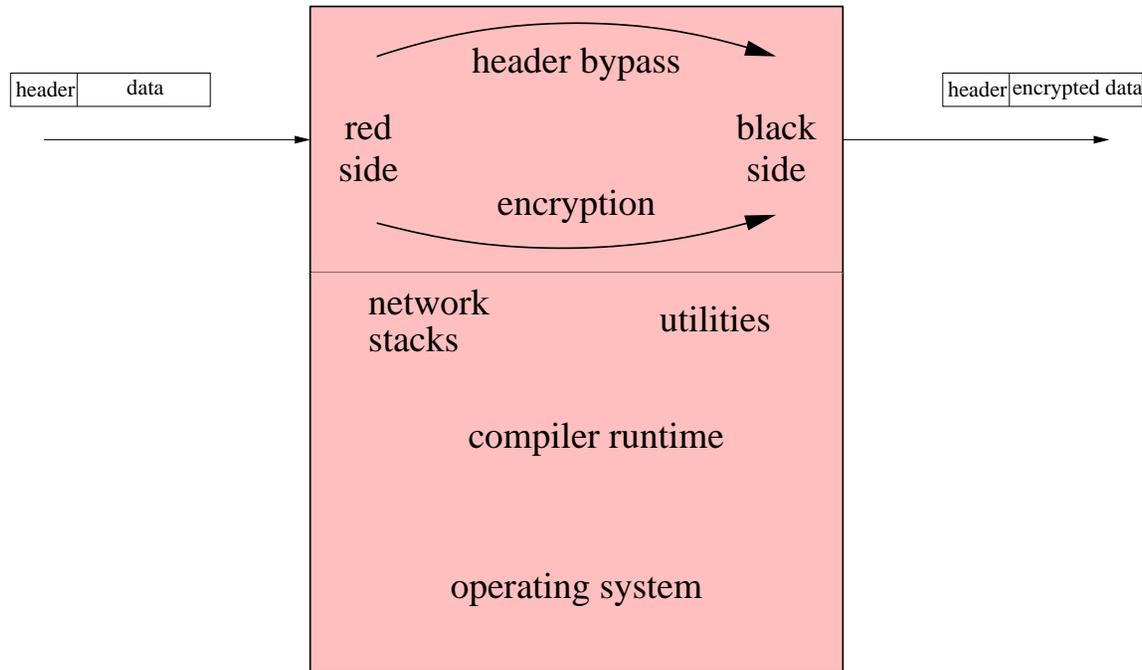
- Construct an architecture so that security assurance does decompose along structural lines
- Two issues in security:
 - Enforce the security policy
 - Manage shared resources securely
- The MILS idea is to handle these separately
- Focus the system architecture on simplifying the argument that policy is enforced correctly
 - Hence policy architecture
- Policy architecture is the interface between the two issues

Policy Architecture

- Intuitively, a boxes and arrows diagram
 - There is a formal model for this
- Boxes encapsulate data, information, control
 - Access only local state, incoming communications
 - i.e., they are state machines
- Arrows are channels for information flow
 - Strictly unidirectional
 - Absence of arrows is often crucial
- Some boxes are trusted to enforce local security policies
- Want the trusted boxes to be as simple as possible
- Decompose the policy architecture to achieve this
- Assume boxes and arrows are free

Crypto Controller Example: Step 1

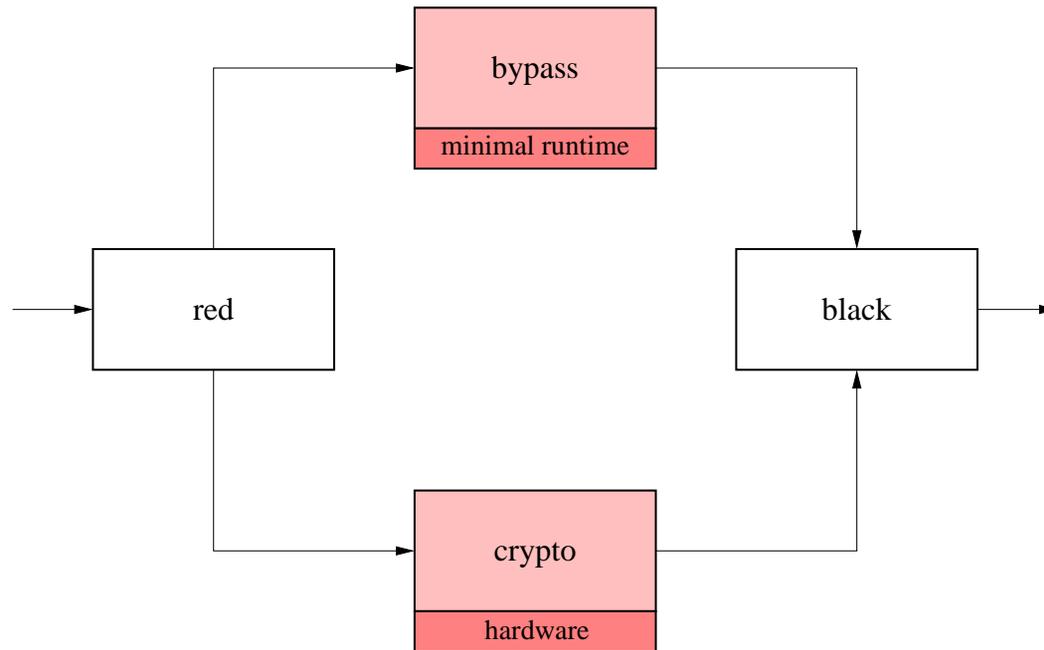
Policy: no plaintext on black network



No architecture, everything trusted

Crypto Controller Example: Step 2

Good policy architecture: fewer things trusted



Local policies (notice these are **intransitive**):

Header bypass: low bandwidth, data looks like headers

Crypto: all output encrypted

Policy Architecture: Compositional Assurance

- Construct assurance for each trusted component **individually**
 - i.e., each component enforces its **local policy**
- Then provide an **argument** that the **local policies**
 - **In the context of the policy architecture**Combine to achieve the **overall system policy**
- **Medium robustness**: this is done informally
- **High robustness**: this is done formally
 - **Compositional verification**

Enforcing Assumptions Of The Policy Architecture

- Primarily **separation**
- Five basic mechanisms available
 - **physical**: separate boxes
 - ★ But even they may need **wrapping**
 - **temporal**: classic periods processing
 - **cryptographic**: encryption and checksums
 - **logical**: verify no interference
 - ★ Only works when you have **all** the code
 - **separation kernel**: runtime enforcement
- Also need **unidirectional arrows**
 - **Data diodes** etc.
- Generally want to combine **separation** with **resource sharing**

Resource Sharing

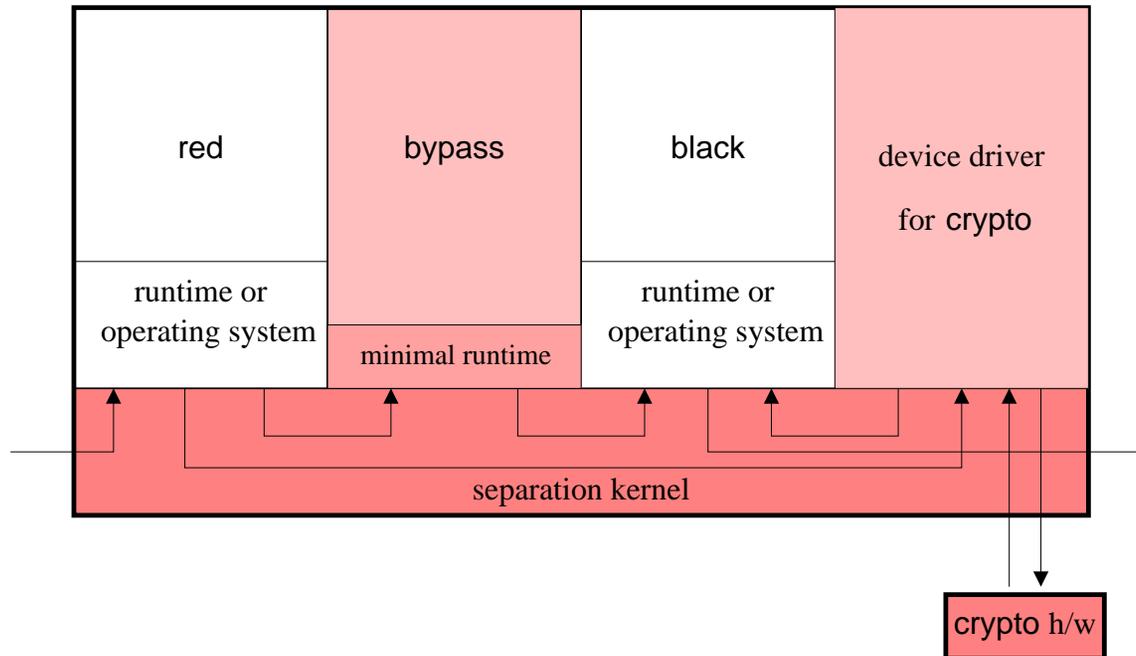
- Next, we need to **implement** the logical components and the communications of **the policy architecture** in an **affordable manner**
- **Allow different components and communications to share resources**
- **Need to be sure the sharing does not violate the policy architecture**
 - Flaws might add new communications paths
 - Might blur the separation between components

Secure Resource Sharing

- For broadly useful classes of resources
 - e.g., file systems, networks, consoles, processors
- Provide implementations that can be shared securely
- Start by defining what it means to partition specific kinds of resource into separate logical components
- Definition in the form of a protection profile (PP)
 - e.g., separation kernel protection profile (SKPP)
 - or network subsystem PP, filesystem PP, etc.
- Then build and evaluate to the appropriate PP

Crypto Controller Example: Step 3

Separation kernel securely partitions the processor resource



The integrity of the policy architecture is preserved

Resource Sharing: Compositional Assurance

- Construct assurance for each resource sharing component **individually**
 - i.e., each component enforces **separation**
- Then provide an **argument** that the individual components
 - Are **additively compositional**
 - e.g., **partitioning(kernel) + partitioning(network)** provides **partitioning(kernel + network)**

And therefore **combine to create the policy architecture**

- **Medium robustness**: this is done informally
- **High robustness**: this is done formally
 - **Compositional verification**

Summary

- MILS (and HAP) are in the mainstream of architectures promoting compositional development and assurance for critical systems
- Ahead in some areas
 - e.g., the policy architecture, COTS cultivation
- Behind in some others
 - e.g., use of elementary control interfaces
 - tool support for assurance
- The challenge ahead is compositional **certification**
- And **regulatory adjustment** to enable this

Thanks

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